

2.7 - Hom Functors, Yoneda's Lemma and Representability

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(Categories)

The material and exposition for this lesson follows an imaginary textbook on Dozzie Abstract Algebra.

A hom functor is a functor $F : \mathbf{C} \rightarrow \mathbf{Set}$ defined in a special way. Though they have a seemingly basic definition, they are actually very important in the future sections of this chapter.

Fix $A \in \text{ob}(\mathbf{C})$, and define F as follows:

1. For each $B \in \text{ob}(\mathbf{C})$, assign FB to the set $\text{hom}(A, B)$;
2. For each $f : B \rightarrow B'$ in \mathbf{C} , define $F(f)$ to be the set map from $\text{hom}(A, B)$ to $\text{hom}(A, B')$ sending $h \rightarrow fh$. [This set map is notated $\text{hom}(A, f)$.]

It is feasibly shown that the data above defines a functor from \mathbf{C} to \mathbf{Set} . This functor is denoted $\text{hom}(A, -)$ and is called a **covariant hom functor**.

Now let $F : \mathbf{C} \rightarrow \mathbf{Set}$ be any functor and let η be a natural transformation $\text{hom}(A, -) \Rightarrow F$. Then for *any* object B in \mathbf{C} , η_B is a set map $\text{hom}(A, B) \rightarrow FB$. In particular, taking A for B , η_A is a set map $\text{hom}(A, A) \rightarrow FA$. Therefore, $a = \eta_A(1_A)$ is some element of FA . We claim that a completely determines the natural transformation η . To see where η_B sends $f : A \rightarrow B$, note the commutativity of

$$\begin{array}{ccc} \text{hom}(A, A) & \xrightarrow{\eta_A} & FA \\ \text{hom}(A, f) \downarrow & & \downarrow F(f) \\ \text{hom}(A, B) & \xrightarrow{\eta_B} & FB \end{array}$$

due to η being a natural transformation. Traveling 1_A along each pair of arrows yields $\eta_B(\text{hom}(A, f)(1_A)) = \eta_B(f1_A) = \eta_B(f)$ and $F(f)(\eta_A(1_A)) = F(f)(a)$. Therefore, $\eta_B(f) = F(f)(a)$, which determines η_B .

It is straightforward to show that any a is possible: if a is an arbitrary element of FA , define $\eta_B : \text{hom}(A, B) \rightarrow FB$ by $\eta_B(f) = F(f)(a)$ for each B . Then

$$\begin{array}{ccc} \text{hom}(A, B) & \xrightarrow{\eta_B} & FB \\ \text{hom}(A, f) \downarrow & & \downarrow F(f) \\ \text{hom}(A, B') & \xrightarrow{\eta_{B'}} & FB' \end{array}$$

is commutative for any $f : B \rightarrow B'$, because one direction yields $F(f)\eta_B$ and the other $\eta_{B'} \text{hom}(A, f)$ and they are the same set map:

$$F(f)(\eta_B(h)) = F(f)(F(h)(a)) = (F(f)F(h))(a) = F(fh)(a) = \eta_{B'}(fh) = \eta_{B'}(\text{hom}(A, f)(h))$$

Therefore, η is a natural transformation $\text{hom}(A, -) \Rightarrow F$. Also, $\eta_A(1_A) = F(1_A)(a) = 1_{FA}(a) = a$. What we have proved is summarized as follows.

LEMMA 2.5 (YONEDA'S LEMMA) *Let \mathbf{C} be a category, $A \in \text{ob}(\mathbf{C})$, $F : \mathbf{C} \rightarrow \mathbf{Set}$ be a functor. Then for every $a \in FA$ there is a unique natural transformation $\eta : \text{hom}(A, -) \Rightarrow F$ such that $a = \eta_A(1_A)$. This natural transformation is given by η_B being the map $f \rightarrow F(f)(a)$ from $\text{hom}(A, B) \rightarrow FB$ for each B .*

Hopefully this sounds like Exercise 3 of Section 1.10 a bit.

Now let $F : \mathbf{V}(S) \rightarrow \mathbf{Set}$ be the forgetful functor. Then let A be the free $\mathcal{V}(S)$ -algebra given by a single symbol a , and let η be the natural transformation $\text{hom}(A, -) \Rightarrow F$ satisfying $\eta_A(1_A) = a$. We have $\eta_B(f) = F(f)(a)$ for each $f : A \rightarrow B$. Since A is the free algebra given by a , then for every $b \in B$ there is a unique homomorphism $f : A \rightarrow B$ such that $f(a) = b$; this implies that η_B is bijective, so that η is actually a natural isomorphism. Hence (A, a) is a special pair for the functor F , which yields the following definition.

DEFINITION

*Let $F : \mathbf{C} \rightarrow \mathbf{Set}$ be a functor. If there exists a pair (A, a) with $A \in \text{ob}(\mathbf{C})$, $a \in FA$ such that the natural transformation $\text{hom}(A, -) \Rightarrow F$ induced by Lemma 2.5 is a natural isomorphism, (A, a) is called a **representative** of the functor F . If functor which has a representative is said to be **representable**.*

EXAMPLES

1. If $F : \mathbf{V}(S) \rightarrow \mathbf{Set}$ is the forgetful functor, we have seen that F is representable with $(F_S(\Omega, \{a\}), a)$ as a representative.

2. Let $F : \mathbf{C} \rightarrow \mathbf{Set}$ be the constant functor onto the one-element set $\{a\}$. Then a representative of F is (I, a) with I an initial object in \mathbf{C} . This is because for every $A \in \mathbf{C}$, $\text{hom}(I, A)$ and FA are both one-element sets, so they come in a unique natural bijection. Hence F is representable if and only if \mathbf{C} has an initial object; in particular, this holds if $\mathbf{C} = \mathbf{V}(S)$.

3. (Quotient algebras) Let A be a $\mathcal{V}(S)$ -algebra and Φ a congruence relation on A . Define $F : \mathbf{V}(S) \rightarrow \mathbf{Set}$ sending B to the subset $\{h \in \text{hom}(A, B) \mid \Phi \subseteq \ker h\}$ of $\text{hom}(A, B)$. Then clearly for $f : B \rightarrow B'$, $\text{hom}(A, f)$ sends elements of FB to elements of FB' , so it can be restricted to a set map $F(f) : FB \rightarrow FB'$. It is easy to see that this data defines a functor.

Now let $\pi : A \rightarrow A/\Phi$ be the canonical epimorphism, then $\pi \in F(A/\Phi)$. We claim that $(A/\Phi, \pi)$ is a representative of F . This is because $\eta_B : \text{hom}(A/\Phi, B) \Rightarrow FB$ sends $h \rightarrow F(h)(\pi) = h\pi$. The injectification theorem (1.10) shows that for all $f : A \rightarrow B$ such that $\Phi \subseteq \ker f$, there is a unique $h : A/\Phi \rightarrow B$ such that $f = h\pi$. This means η_B is actually bijective, so η is a natural isomorphism.

4. (Colimits) Let $D : \mathbf{J} \rightarrow \mathbf{C}$ be a fixed diagram of type \mathbf{J} in \mathbf{C} . Define $F : \mathbf{C} \rightarrow \mathbf{Set}$ sending each object B of \mathbf{C} to the set of cocones $(B, \{\eta_\alpha\})$ from D to B . If $f : B \rightarrow B'$ is a morphism in \mathbf{C} and $(B, \{\eta_\alpha\})$ is a cocone from D , clearly $(B', \{f\eta_\alpha\})$ is also a cocone; this induces a set map $F(f) : FB \rightarrow FB'$. It is straightforward to show that F is a functor.

Now let $(L, \{\eta_\alpha\})$ be a colimit of D . We claim that L , along with this cocone, is a representative of F . The natural transformation $\eta : \text{hom}(L, -) \Rightarrow F$

assigns each $\eta_B, B \in \text{ob}(\mathbf{C})$ the map $\theta \rightarrow F(\theta)(L, \{\eta_\alpha\}) = (B, \{\theta\eta_\alpha\})$ from $\text{hom}(L, B) \rightarrow FB$. Recall that the virtue of being a colimit is that for any cocone $(B, \{\zeta_\alpha\})$, there is a unique $\theta : L \rightarrow B$ such that $\zeta_\alpha = \theta\eta_\alpha$ for every α ; that is, $(B, \{\zeta_\alpha\}) = \eta_B(\theta)$. Therefore η_B is bijective.

Conversely, every representative of F is a colimit of D ; see Exercise 5 for proof.

We now fix $A \in \text{ob}(\mathbf{C})$ and define a contravariant functor as follows.

1. For each $B \in \text{ob}(\mathbf{C})$, assign FB to the set $\text{hom}(B, A)$.
2. For each $f : B \rightarrow B'$ in \mathbf{C} , define $F(f)$ to be the map $\text{hom}(B', A) \rightarrow \text{hom}(B, A)$ sending $h \rightarrow hf$. [This set map is notated $\text{hom}(f, A)$; be wary that B and B' swap places.]

This functor is denoted $\text{hom}(-, A)$ and is called a **contravariant hom functor**. Another way to view this is that $\text{hom}_{\mathbf{C}}(-, A) = \text{hom}_{\mathbf{C}^{\text{op}}}(A, -)$ which is a covariant functor $\mathbf{C}^{\text{op}} \rightarrow \mathbf{Set}$. Yoneda's Lemma then dualizes to $\text{hom}(-, A)$ as follows:

LEMMA 2.5 (CONTINUED) *Let $F : \mathbf{C} \rightarrow \mathbf{Set}$ be a contravariant functor, $A \in \text{ob}(\mathbf{C})$. Then for every $a \in FA$ there is a unique natural transformation $\eta : \text{hom}(-, A) \Rightarrow F$ such that $a = \eta_A(1_A)$. This natural transformation is given by η_B being the map $f \rightarrow F(f)(a)$ from $\text{hom}(B, A) \rightarrow FB$ for each B .*

This immediately follows from the first one when \mathbf{C}^{op} is used. When the pair (A, a) is special enough for η to be a natural isomorphism, F is said to be **representable** with (A, a) as a representative.

EXAMPLES

1. Let $F : \mathbf{C} \rightarrow \mathbf{Set}$ be the contravariant constant functor onto the one-element set $\{a\}$. Then a representative of F is (T, a) with T a terminal object in \mathbf{C} . Hence F is representable if and only if \mathbf{C} has a terminal object.

2. (Subalgebras) Let A be a $\mathcal{V}(S)$ -algebra and A' a subalgebra on A . Define $F : \mathbf{V}(S) \rightarrow \mathbf{Set}$ sending B to the subset $\{h \in \text{hom}(B, A) \mid h(B) \subseteq A'\}$ of $\text{hom}(B, A)$. Then clearly for $f : B \rightarrow B'$, $\text{hom}(f, A)$ sends elements of FB' to elements of FB , so it can be restricted to a set map $F(f) : FB' \rightarrow FB$. This data defines a contravariant functor.

Now let $\iota : A' \rightarrow A$ be the canonical monomorphism, then $\iota \in FA'$. We claim that (A', ι) is a representative of F . This is because $\eta_B : \text{hom}(B, A') \Rightarrow FB$ sends $h \rightarrow F(h)(\iota) = \iota h$. The bijectivity of each η_B follows from the surjection theorem (1.5).

3. (Limits) Let $D : \mathbf{J} \rightarrow \mathbf{C}$ be a fixed diagram of type \mathbf{J} in \mathbf{C} . Define $F : \mathbf{C} \rightarrow \mathbf{Set}$ sending each object B of \mathbf{C} to the set of cones $(B, \{\eta_\alpha\})$ from B to D . Then for $f : B' \rightarrow B$ a morphism in \mathbf{C} and $(B, \{\eta_\alpha\})$ a cone to D , so is $(B', \{\eta_\alpha f\})$, inducing a set map $F(f) : FB \rightarrow FB'$. This defines a contravariant functor F . The reasoning above shows that a representative of F is a limit of the diagram D .

4. Let \mathcal{P} be the contravariant functor of Example 3 of Section 3. Now let \mathbb{Z}_2 be the set $\{0, 1\}$, then $\{1\}$ is a subset of \mathbb{Z}_2 , hence is in $\mathcal{P}(\mathbb{Z}_2)$. We claim that \mathcal{P} is representable with $(\mathbb{Z}_2, \{1\})$ as a representative. The natural transformation $\eta : \text{hom}(-, \mathbb{Z}_2) \Rightarrow \mathcal{P}$ sends a set X to the map $\eta_X : \text{hom}(X, \mathbb{Z}_2) \Rightarrow \mathcal{P}(X)$ sending $f \rightarrow F(f)(\{1\}) = f^{-1}(\{1\}) = \{x \in X \mid f(x) = 1\}$. Since for every subset X' of X , the *characteristic function* $f : X \rightarrow \mathbb{Z}_2$ [$f(x) = 1$ for $x \in X'$, 0 for $x \in X - X'$] is the unique element of $\text{hom}(X, \mathbb{Z}_2)$ sent to X' by η_X , η_X is bijective. Therefore, η is a natural isomorphism.

The hom functors have a huge use in the next section.

EXERCISES

1. Use Yoneda's Lemma to show that for every $h : A' \rightarrow A$ in \mathbf{C} , there is a unique natural transformation $\eta : \text{hom}(A, -) \Rightarrow \text{hom}(A', -)$ such that $h = \eta_A(1_A)$.
2. Let \mathbf{C} be an arbitrary category, and define $H : \mathbf{C}^{\text{op}} \times \mathbf{C} \rightarrow \mathbf{Set}$ as follows. For each pair (A, B) of objects of \mathbf{C} , $H(A, B) = \text{hom}(A, B)$. To define $H(f)$ for $f : (A, B) \rightarrow (A', B')$ in $\mathbf{C}^{\text{op}} \times \mathbf{C}$, note that f is a pair of morphisms $f_1 : A' \rightarrow A, f_2 : B \rightarrow B'$ of \mathbf{C} . Assign $H(f)$ the set map $h \rightarrow f_2 h f_1$ from $\text{hom}(A, B)$ to $\text{hom}(A', B')$. Verify that this is a functor. It is called the **two-variable hom functor**.
3. Show that $\text{hom}(A, -)$ is continuous. [See Exercise 5 of Section 6.] Conclude that every [covariant] representable functor is continuous.
4. Let X be a set, $A \in \text{ob}(\mathbf{C})$, and $A' = \coprod_{x \in X} A$ be a copower of A with injections $i_x : A \rightarrow A', x \in X$. Then define $u : X \rightarrow \text{hom}(A, A')$ sending $x \rightarrow i_x$. Show that (A', u) is a universal from X to $\text{hom}(A, -)$.
5. A representative of a functor F is precisely a universal from the one-element set $\{o\}$ to F . Conclude that if (A, a) and (A', a') are two representatives of F , there is a unique isomorphism $\sigma : A \rightarrow A'$ such that $F(\sigma)(a) = a'$.
6. Let $f : A' \rightarrow A, g : B \rightarrow B'$ in \mathbf{C} . Explain why

$$\begin{array}{ccc} \text{hom}(A, B) & \xrightarrow{\text{hom}(f, B)} & \text{hom}(A', B) \\ \text{hom}(A, g) \downarrow & & \downarrow \text{hom}(A', g) \\ \text{hom}(A, B') & \xrightarrow{\text{hom}(f, B')} & \text{hom}(A', B') \end{array}$$

is a commutative diagram. Use this to show that $F : \mathbf{C} \rightarrow \mathbf{Set}^{\mathbf{C}}$ is a contravariant functor when one defines $FA = \text{hom}(A, -)$ and for $f : A' \rightarrow A$ in \mathbf{C} , $F(f) : \text{hom}(A, -) \Rightarrow \text{hom}(A', -)$ is the natural transformation given by $\eta_B = \text{hom}(f, B) : \text{hom}(A, B) \rightarrow \text{hom}(A', B)$.